



# MA014G

## Algebra och Diskret Matematik A

### Svar på uppgifter till Block 6

Referenser utan parenteser är till [J] edition 5, referenser i ()-parenteser är till [J] edition 4, och referenser i []-parenteser är till [J] edition 6.

## Uppgifter i läsanvisningen

### Uppgift B6.1 - B6.4

[J] Section 6.1 (6.1) [8.1]:

6 (2) [6].

The edge  $ab$  must be included, so we start with it. Now we have a choice:  $bc$  or  $bd$ . If we take  $bc$  we are left with the edge  $bd$ , which we have to take at some point. When we do take it, we will find ourselves stuck at  $b$  with no edge remaining by which to leave  $b$ . So we cannot take  $bc$ . If instead we take  $bd$ , we find we have the same problem, this time with  $bc$  remaining.

7 (3) [7].

We will have exactly the same problem as in exercise 6 (2) [6]. When we go through  $b$  we use up two edges, so we can go through it twice, but when we visit it the third time, we use the last edge and cannot leave again.

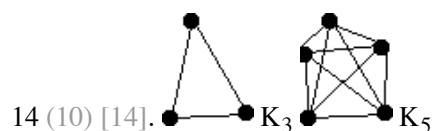
9 (5) [9].  $abdefcebc$

10 (6) [10].  $abcfihfehgecgdebda$

11 (7) [11]. Not simple as it contains both a loop ( $e_5$ ) and parallel edges ( $e_1$  and  $e_6$ ).  $e_1$  is incident on  $v_1$  and  $v_2$ .

12 (8) [12]. Simple graph.  $e_1$  is incident on  $v_2$  and  $v_4$ .

13 (9) [13]. Simple graph. There is no  $e_1$ .



14 (10) [14].  $n(n-1)/2$ .

15 (11) [15].  $n(n-1)/2$ .

16 (12) [16].  $K_{3,3}$  - see solution to exercise 20 (16) [24] in [J].

17 (14) [18]. Bipartite.  $V_1 = \{v_1, v_6, v_3, v_4, v_8, v_{10}, v_9\}$   $V_2 = \{v_5, v_2, v_7\}$

18 (15) [19]. Bipartite.  $V_1 = \{\text{Gre, Buf, Dou, Sho, Mud}\}$   $V_2 = \{\text{She, Wor, Cas, Lan, Gil}\}$

19 (17) [21]. Not bipartite as it contains a triangle,  $v_1 v_2 v_3 v_1$ .

20 (18) [22]. See 21 (17) [21].

21 (21) [25].  $mn$

### Uppgift B6.5

1,2,5,6,2,1 är en stig från 1 till 1, men den är inte enkel.

1,2,5,7 är inte en stig, ty det finns ingen kant (5,7) i grafen.

## Uppgift B6.6 - B6.7

[J] Section 6.2 (6.2) [8.2]:

2 (2) [2]. simple path

3 (3) [3]. none

5 (5) [5]. none

6 (6) [6]. cycle

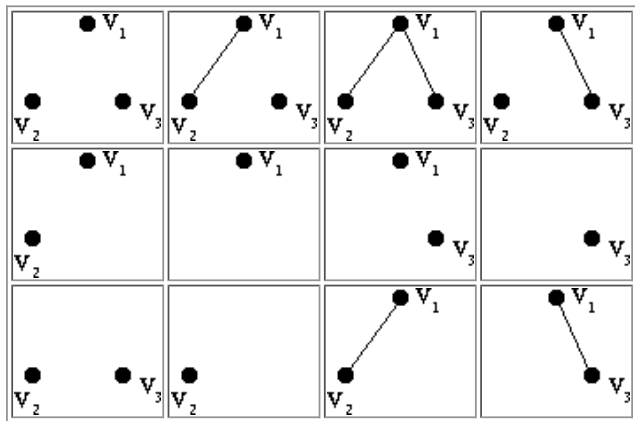
8 (8) [8]. simple path

9 (9) [9]. simple path

12 (12) [12].



26 (26) [26].



29 (29) [29]. Has an Eulerian cycle as all degrees are even.

30 (30) [30]. No Eulerian cycle as  $v_3$  has odd degree.

32 (32) [32]. abdcfbfgjfejhcdieih

33 (33) [33]. abcdehfgijhkfegfdca

35 (35) [35]. when  $n$  is odd

36 (36) [36]. both  $m$  and  $n$  must be even

## Uppgift B6.8

[J] Section 6.3 (6.3) [8.3]:

2 (2) [2]. abcdefnpmlkjoihga

4 (4) [4]. This graph is bipartite, with an uneven partition of the vertices so it is not Hamiltonian.

5 (5) [5]. The edges  $gi$ ,  $ij$ ,  $fg$ ,  $gh$  and  $jg$  would all have to be included, which would lead to  $g$  being used too many times.

7 (7) [7]. abcgImrqpkjfeinotshda

8 (8) [8].

We can only use 2 edges at c, 2 at e and 2 at f so we don't use 2 edges at c, 1 edge at d and 3 edges at e leaving a total of 6 edges which is too few for a hamilton cycle.

10 (10) [10].  $K_3$

11 (11) [11].  $K_5$

## Uppgift B6.9

[J] Section 6.4 (6.4) [8.4]:

2 (2) [2]. 11 abcg

3 (3) [3]. 10 abcdz

5 (5) [5]. 10 hfed

## Uppgift B6.10

[J] Section 6.5 (6.5) [8.5]:

2 (2) [2].

$$\begin{array}{c}
 \begin{array}{ccccccc}
 & a & b & c & d & e & f & g
 \end{array} \\
 \begin{array}{c}
 a \\
 b \\
 c \\
 d \\
 e \\
 f \\
 g
 \end{array}
 \begin{pmatrix}
 0 & 2 & 0 & 0 & 0 & 0 & 0 \\
 2 & 0 & 1 & 0 & 1 & 0 & 0 \\
 0 & 1 & 2 & 0 & 1 & 0 & 0 \\
 0 & 0 & 0 & 2 & 1 & 0 & 1 \\
 0 & 1 & 1 & 1 & 0 & 1 & 0 \\
 0 & 0 & 0 & 0 & 1 & 0 & 1 \\
 0 & 0 & 0 & 1 & 0 & 1 & 0
 \end{pmatrix}
 \end{array}$$

8 (8) [8].

$$\begin{array}{c}
 \begin{array}{cccccccccc}
 & x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & x_8 & x_9 & x_{10}
 \end{array} \\
 \begin{array}{c}
 a \\
 b \\
 c \\
 d \\
 e \\
 f \\
 g
 \end{array}
 \begin{pmatrix}
 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\
 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0
 \end{pmatrix}
 \end{array}$$

## Uppgift B6.11

### [J] Section 6.6 (6.6) [8.6]:

2 (2) [2].

Isomorphic, as there is a bijection  $f$  from the vertices of  $G_1$  to the vertices of  $G_2$  which preserves the edges.  $f$  is defined by  $f(a)=a'$ ,  $f(b)=c'$ ,  $f(c)=e'$ ,  $f(d)=g'$ ,  $f(e)=b'$ ,  $f(f)=d'$ ,  $f(g)=f'$ .

3 (3) [3].

These are not isomorphic. Choose ONE of the following to show this:

- $G_1$  contains triangles (such as  $abca$ ) whereas  $G_2$  contains no triangles.
- The vertices of  $G_2$  cannot be coloured with 2 colours so that no two adjacent vertices are coloured with the same colour.  $G_2$  can be coloured with 2 colours.
- $G_2$  is bipartite whereas  $G_1$  is not.
- The complement of  $G_1$  consists of one component whereas the complement of  $G_2$  consists of two.

This list is probably not complete, try to think of some other structural differences yourself.

**note that** giving a failed attempt to pair up the vertices is **never enough** to show that two graphs are non-isomorphic unless you have tried **every** possible way to pair them up - that is, in this case all  $6! = 720$  ways.

5 (5) [5].

Isomorphic, as there is a bijection  $f$  from the vertices of  $G_1$  to the vertices of  $G_2$  which preserves the edges.  $f$  is defined by  $f(a)=b'$ ,  $f(b)=d'$ ,  $f(c)=e'$ ,  $f(d)=a'$ ,  $f(e)=c'$ .

6 (6) [6].

Here is a description of how I solved this problem. Note that such a description would not usually be included in an answer given to a question of this kind, one would just give the 'answer' below. First, I wrote the degree of each vertex next to the vertex in the graph, and then compared the vertices of each degree. There are two vertices of degree 5 in each graph, and as  $G_2$  is symmetrical, it doesn't matter which way we pair them up. Take  $f(l)=l'$  and  $f(k)=k'$ . Now there are two vertices of degree 4 adjacent to both  $l$  and  $l'$ . Try pairing those (might be wrong - we'll see):  $f(a)=a'$ ,  $f(h)=d'$ . Now  $k$  is adjacent to two vertices of degree 4, namely,  $e$  and  $d$ , where  $d$  is adjacent to  $a$ . Similarly  $k'$  is adjacent to  $c'$  and  $b'$  where  $b'$  is adjacent to  $f(a)=a'$ . To preserve the edges, we pair  $d$  and  $b'$  so  $f(d)=b'$ , leaving  $f(e)=c'$ . Now I look for vertices which are adjacent to those I have already paired up.  $g$  is adjacent to both  $l$  and  $h$ , and  $h'$  is adjacent to both  $f(l)=l'$  and  $f(h)=d'$  so  $f(g)=h'$ .  $f$  is adjacent to  $g, k, e$  and  $g'$  is adjacent to  $f(g)=h'$ ,  $f(k)=k'$  and  $f(e)=c'$  so  $f(f)=g'$ . Only one of the neighbours of  $a$  is not paired and that is  $b$ . Similarly, of the neighbours of  $f(a)=a'$ , only  $e'$  is unpaired so  $f(b)=e'$ .  $c$  is adjacent to both  $b$  and  $d$  and  $f'$  is adjacent to both  $f(b)=e'$  and  $f(d)=b'$  so  $f(c)=f'$ . We are left with just  $i$  and  $j$  and  $i'$  and  $j'$ . By comparing neighbours, we see  $f(i)=i'$  and  $f(j)=j'$ . Now for the 'answer' to the question:

$G_1$  and  $G_2$  are isomorphic because there is a bijection  $f$  mapping the vertices of  $G_1$  to  $G_2$  so that there is an edge  $uv$  in  $G_1$  if and only if the corresponding edge  $f(u)f(v)$  is present in  $G_2$ . The function  $f$  is defined as follows:  $f(a)=a'$ ,  $f(b)=e'$ ,  $f(c)=f'$ ,  $f(d)=b'$ ,  $f(e)=c'$ ,  $f(f)=g'$ ,  $f(g)=h'$ ,  $f(h)=d'$ ,  $f(i)=i'$ ,  $f(j)=j'$ ,  $f(k)=k'$ ,  $f(l)=l'$ .

8 (8) [8].

Hint: look at the vertices of degree 3.

9 (9) [9].

First I looked at the degrees. It didn't help. Next I looked at the complements. Did not look too nice, so I decided to come back to that later if I could not find anything else. Then I noticed that in both graphs the only vertex of degree 4 is contained in two triangles in  $G_1$  and just one triangle in  $G_2$ .

$G_1$  and  $G_2$  are not isomorphic because (chose any ONE of the following)

- the only vertex  $e$  of degree 4 in  $G_1$  is contained in two triangles  $abe$  and  $efg$ , whereas the only vertex  $f'$  of degree 4 in  $G_2$  is contained in just one triangle  $f'g'e'$ .
- if we remove the vertices of degree 4 from  $G_1$  and from  $G_2$ , one of the remaining graphs can be coloured with 2 colours whereas the other can not.
- $G_1$  contains two triangles whereas  $G_2$  contains only one.

## Uppgift B6.12

### [J] Section 7.1 (7.1) [9.1]:

9 (2) [9].  
Answer 4

10 (4) [10].  
Answer 5

## Uppgift B6.13

### [J] Section 7.2 (7.2) [9.2]:

2 (2) [2]. Aphrodite, Uranus.

3 (3) [3]. Aphrodite, Kronos, Atlas, Prometheus.

5 (5) [5]. Zeus, Poseidon, Hades.

6 (6) [6]. One edge drawn vertically with Aphrodite as root and Eros as leaf.

24 (24) [24]. No such graph exists. A terminal node (leaf) has degree one.

26 (26) [26]. The root has degree three and each child of the root is adjacent to two leaves.

## Uppgift B6.14

### [J] Uppgift 7.3.9 (7.3.9) [9.3.9]:

The edges in the spanning tree are (a,c), (b,c), (c,d), (d,f), (f,e), (f,g), (f,h), (f,i).

## Uppgift B6.15

### [J] Section 7.4 (7.4) [9.4]:

2 (2) [2].  
The edges in the minimal spanning tree are (1,2), (2,3), (3,6), (6,9), (2,5), (5,8), (8,7), (7,4).

3 (3) [3].  
The edges in the minimal spanning tree are (4,1), (1,2), (2,5), (5,6), (6,3).

5 (5) [5].  
The edges in the minimal spanning tree are  
(1,2), (2,5), (2,7), (3,7), (4,8), (5,6), (5,9), (6,15), (7,11), (8,12), (9,10), (10,13), (11,14), (12,16), (15,16).

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